

THALES

Building a future we can all trust

A Constraint Programming Approach to Ship Refit Project Scheduling

Raphaël Boudreault, *Thales Digital Solutions (Québec, Canada)*

Vanessa Simard, *NQB.ai (Québec, Canada)*

Daniel Lafond, *Thales Digital Solutions (Québec, Canada)*

Claude-Guy Quimper, *Université Laval (Québec, Canada)*

CP 2022

August 2, 2022



What is a ship refit?

- Important shipyard event where all ship's activities are **suspended**
- Objective is to restore, customize, modify or modernize part of a ship
- Made of **several hundred (or thousand) tasks**
- Can span over **several weeks, months (or over a year)**
- **Longer refit = higher costs**
- Time window must be planned **years in advance**
- When exceeded, the dock must be cleared



Ship refit planning

- **Complex and tedious**
- Initial planning (free of conflicts) can take up to 120 days
- Day-to-day re-planning is **difficult and time-consuming**
- Typical software (Primavera P6, Microsoft Project) have **limited optimization capabilities** (not exact, only resources leveling, etc.)

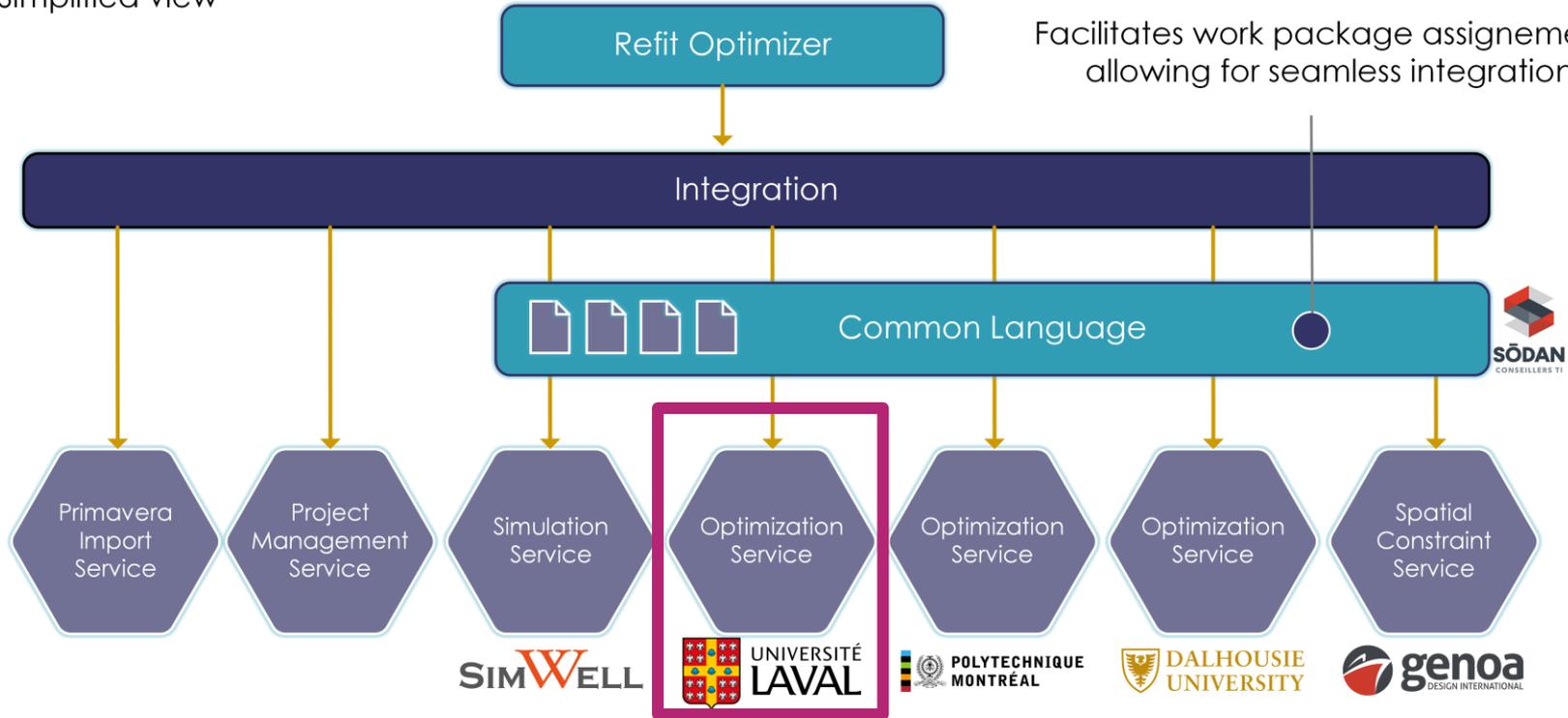


Refit Optimizer

- Prototype solution for **multi-objective optimization** in the ship refit domain
- **Generic** architecture for other scheduling contexts
- Key motivation: Challenges identified in the **Arctic and Offshore Patrol Ship and Joint Support Ship In-Service Support (AJISS) program** with the Royal Canadian Navy

Simplified view

Facilitates work package assignments allowing for seamless integration



Operational and deployed on a secured cloud platform (Thales TrustNest)

Elements to consider

- Planning **horizon**
- Planning **granularity** (days or hours)
- Tasks depend on **capacity-limited resources** (human/material)
- Maximum number of workers simultaneously in some **work areas**
- **Precedence relationships** between tasks
- **Date constraints** (e.g. milestones)
- Some tasks must be **idle during weekends**
- Some tasks *can* be performed in **overtime**

Objectives

1. Minimize the refit total duration (**makespan**)
2. Minimize the costs associated with overtime labor (**overtime**)
3. Minimize the risk, planning the overtime as early as possible (**robustness**)

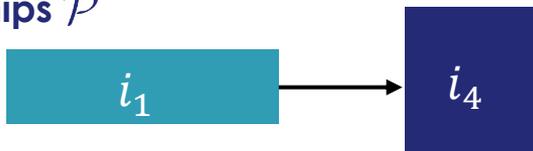
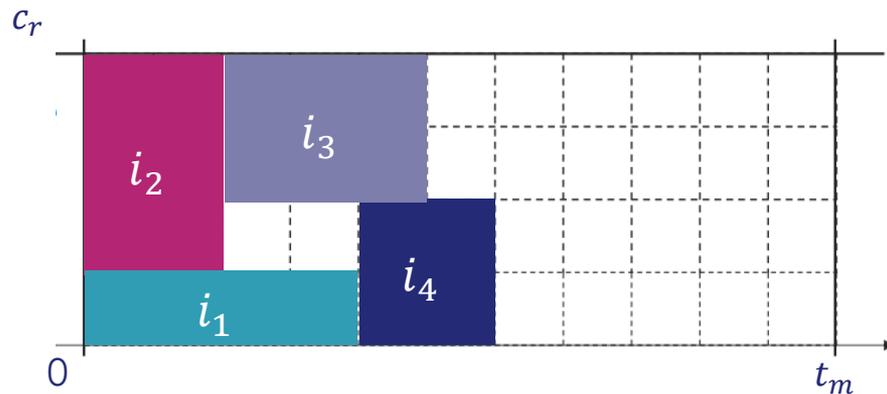
Instance	Horizon	#Tasks (#O)	Task duration	#Precedence relations	#Resources (#WA)
Artificial	day-yacht21	21 (20)	1-3 days	32	9 (2)
	hour-yacht21	21 (20)	1-8 hours	32	9 (2)
	generic136	136 (136*)	1-20 days	99	9 (4)
	software138	138 (138*)	1-10 days	341	8 (0)
Realistic	navy253	253 (253*)	1-8 hours	246	92 (87)
	cruise510	510 (464*)	1-15 days	550	32 (24)
	navy830	830 (830*)	1-200 hours	816	146 (128)

Can be performed in overtime
*Must be idle during weekends

Work areas

Standard definition (Pritsker et al., 1969)

- Set of **tasks** \mathcal{I}
- **Timeline** $\mathcal{T} = \{0, 1, \dots, t_m\}$, **horizon** t_m
- Set of **resources** \mathcal{R}
- Task $i \in \mathcal{I}$ **requires** $h_{i,r}$ of resource $r \in \mathcal{R}$, for its whole duration
- Each resource $r \in \mathcal{R}$:
 - **Capacity** c_r
 - **Renewable** (fully available at all time)
 - **Cumulative** (more than one task can use a resource at a time)
- Set of **precedence relationships** \mathcal{P}



- Objective: Find a schedule with the **minimal makespan**

- **Significant efforts** in the CP community to solve scheduling problems with resources
- **CUMULATIVE global constraint** (Aggoun & Beldiceannu, 1993)

$$\sum_{\substack{i \in \mathcal{I}: \\ S_i \leq t < S_i + D_i}} h_{i,r} \leq c_r \quad \forall t \in \mathcal{T}. \quad \longleftrightarrow$$

Usage of a resource is at most its capacity for each time point in the timeline

- Many filtering rules developed and improved
 - Time-Tabling
 - Time-Table Edge Finding (TTEF)
 - Energetic Reasoning...
- **Important progress towards solving large-scale RCPSP** (Schutt *et al.*, 2011, 2013)
- **Lazy clause generation** (Ohrimenko *et al.*, 2009)
 - Hybrid between CP and **SAT solvers**
 - Filtered values recorded with explanations as SAT clauses
 - On a failure, learns a *nogood*
 - Solvers: **Chuffed**, OR-Tools (Google), CP Optimizer (IBM)
 - SAT-based branching heuristic: *Variable State Independent Decaying Sum* (**VSIDS**)

Planning granularity in days

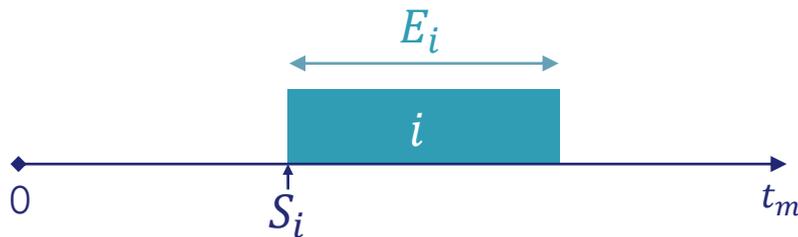
Additional parameters

- $s_i^U, s_i^L, e_i^U, e_i^L$, bounds on start/end times implied by **date constraints**
- p_i , processing time (**task duration without overtime**)
- w_r^S, w_r^O , **daily standard/overtime usage cost** of resource r ($w_r^S \leq w_r^O$)
- A working day schedule:
 

The diagram shows a horizontal timeline with three points: d_S , d_O , and d_E . A blue double-headed arrow labeled "Standard" spans from d_S to d_O . A red double-headed arrow labeled "Overtime" spans from d_O to d_E .
- Set of tasks that can be planned with overtime $\mathcal{I}^* \subseteq \mathcal{I}$

Decision variables

- For each task $i \in \mathcal{I}$
 - **Starting time** $s_i \in [s_i^L, s_i^U]$
 - **Elapsed time** $E_i \in \mathcal{T}$



Constraints

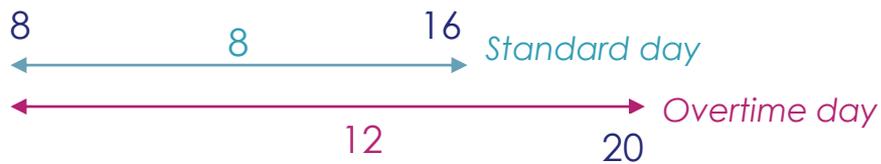
$$\text{CUMULATIVE}([S_i \mid i \in \mathcal{I}], [E_i \mid i \in \mathcal{I}], [h_{i,r} \mid i \in \mathcal{I}], c_r) \quad \forall r \in \mathcal{R}$$

$$e_i^L \leq S_i + E_i \leq e_i^U \quad \forall i \in \mathcal{I}$$

$$S_i + E_i + l \leq S_j \quad \forall (i, j, l) \in \mathcal{P}$$

$$E_i = p_i \quad \forall i \in \mathcal{I} \setminus \mathcal{I}^*$$

$$\left\lceil \frac{(d^O - d^S) p_i}{d^E - d^S} \right\rceil \leq E_i \leq p_i \quad \forall i \in \mathcal{I}^*$$



$$\left\lceil \frac{8 * 3}{12} \right\rceil \leq E_i \leq 3$$

Objectives

1. Makespan

$$\min \max_{i \in \mathcal{I}} (S_i + E_i)$$

$$\mathcal{I}^* = \emptyset$$

2. Overtime

$$\min \sum_{i \in \mathcal{I}^*} \overbrace{(p_i - E_i)}^{\text{Overtime days}} \left(\sum_{r \in \mathcal{R}} h_{i,r} \overbrace{(w_r^O - w_r^S)}^{\text{Daily overtime cost/resource}} \right)$$

3. Robustness

$$\min \sum_{i \in \mathcal{I}^*} (p_i - E_i) S_i$$

- More types of **precedence constraints**

$$X_i \pm l \leq Y_j$$

- Suspension of some tasks during **weekends**
 - Additional variables N_i , **non-working (idle) time points**
 - Included in the elapsed time with specific constraints
- Support of **planning granularity in hours**
 - Additional variable O_i , **overtime time points**
 - Constraints for relation with N_i , which includes *nights*
 - Elapsed time is **replaced** by $p_i + N_i$

BASELINE strategy

> Makespan

S_i with smallest value in domain, assigned to that value

Focus: Scheduling as early as possible

> Overtime and robustness

1. Choose i such that S_i has smallest value in domain
2. Assign smallest value to S_i
3. Assign greatest value to E_i

Focus: Scheduling as early as possible with as few overtime as possible

Formulated as a priority search in MiniZinc

SBPS strategy

- Uses a simple and efficient **value selection heuristic**
 - *Best-Solution* (Vion and Piechowiak, 2017)
 - ***Solution-Based Phase Saving*** (SBPS) (Demirović *et al.*, 2018)

If b is the value of X in the **current best solution**, when branching on X :

- If b is in domain of X , **choose b**
- Else, use a **fallback heuristic**

- Combined with a restart strategy and a dynamic variable selection heuristic, effectively mimics a *Large Neighborhood Search* (LNS), **without loss of exactness**
- We use *BASELINE* until a first solution
- Then, use SBPS with **conflict activity (VSIDS)** variable selection and *BASELINE* as fallback

Setup

- Modeled with *MiniZinc*
- SBPS scheme implemented in *Chuffed* CP solver, that we used
- CUMULATIVE set to use TTEF checking and filtering
- Timeout: 4 hours
- Constant restart strategy of 100 failures

Experiments

- Each instance, each objective, each strategy
- **Overtime/Robustness:** restricted horizon between 2-30% of the best known makespan
 - Not *generic 136*, due to special structure

■ **Table 3** Results on the benchmark instances when considering the **makespan** objective.

Instance	<i>BASELINE</i>		<i>SBPS</i>		Time (s) improv.
	Objective	Time (s)	Objective	Time (s)	
day-yacht21	28 days	0.2*	28 days	0.2*	0.2
hour-yacht21	78 hours	0.4*	78 hours	0.4*	0.4
generic136	178 days	0.7*	178 days	0.7*	0.7
software138	144 days	1.4	119 days	41.6	1.1
navy253	389 hours	4.2	389 hours	3.7	3.7
cruise510	228 days	14.7	227 days	785.7	229.3
navy830	5216 hours	18.7	5144 hours	199.7	18.2

Best makespan reduced by 5% on average

■ **Table 4** Results on the benchmark instances when considering the **overtime** objective.

Instance	BASELINE		SBPS		Time (s) improv.
	Objective	Time (s)	Objective	Time (s)	
day-yacht21	1560	0.3*	1560	0.3*	0.3
hour-yacht21	485	0.4*	485	0.4*	0.4
software138	5600	14 359.6	2600	153.4	34.3
navy253	70	4.2	66	5.0	4.0
cruise510	26 000	11.7	15 760	7555.3	5.8
navy830	227	25.2	36	276.5	26.6

*Best cost reduced by
48% on average*

■ **Table 5** Results on the benchmark instances when considering the **robustness** objective.

Instance	BASELINE		SBPS		Time (s) improv.
	Objective	Time (s)	Objective	Time (s)	
day-yacht21	47	0.3*	47	0.3*	0.3
hour-yacht21	192	0.4*	192	0.4*	0.4
software138	900	13 571.8	258	320.2	15.3
navy253	10 686	5057.6	3480	1411.9	6.7
cruise510	4870	13 022.5	842	1321.7	14.2
navy830	146 794	11 208.9	9076	13 863.4	41.1

*Best value reduced by
79% on average*

➤ Solving time restrictions

- Obtain “good” solutions under 15 min. for < 100 tasks, under 4 hours for > 500 tasks
- In comparison, **up to 4 hours** to manually “optimize” *day-yacht21*

➤ Anonymity

- Estimated workforce costs changed to abstract values

➤ Explainability of results

- Input data format, parameter selection, etc.
- Focus on results interpretation and solution selection
- **Unsatisfiability** during initial planning of real projects

Contributions

- Introduced a CP approach for the ship refit planning problem
- **Successfully tested** on seven industrial instances
 - Detailed complexity analysis with RCPSP metrics in the paper
 - Three objective functions (**makespan, overtime, robustness**)
 - Used SBPS value selection to speed-up the search
 - Better solutions found **significantly faster** than baseline

Next steps

- Complex geospatial constraints and visualization (Lafond *et al.*, 2021)
- Experiments with *Mixed-Integer Programming* model
- Consider task priority levels
- Further explore simulations for robustness assessment
- **Maintenance Optimizer**: long-term planning of preventive maintenance over work periods